

29/05/19

# NORMALISATION



= Normalisation - process of analysing and correcting the table structure so as to minimise data redundancy (complication) and minimise the insertion/deletion/update anomalies.

Proj. Num	Proj. Name	Emp. num	Emp. name	Job class	change Hours	Hours worked	10/1/19
15	Research	105	Aman	elec. Engg.	10		
?	-	107	Abha	BSS Analyst	15		
-	-	110	Raman	DB. Anlgner	17		
-	-	-	Konik	EE			
16	Development	106	<del>etc E</del> Bob	etc E	18		
-	-	104	<del>Data</del> Bob	Data Sca	20		
-	-	103	DBA Dylan	DBA	100		
-	-	115	<del>Sys Analyst</del>	sync. Analyst	1500		
17	Manufacture	113	Amit		160		
-	-	117	Meeraj		115		



## Functional Dependencies

The attribute 'B' is fully functional dependent on attribute 'A' if each value of 'A' determines one and only one value of 'B'. eg:-

Proj. Num  $\rightarrow$  Proj. Name

Project Proj. Num functionally determines Proj. Name which means for every project no there is a name associated with it

## (2) Fully functional dependencies

If attribute 'B' is functionally dependent on a composite key A (composite - multiple attribute) but not on any subset of that composite key then the attribute B

is fully functionally dependent on A. eg:  
~~proj-no.~~ proj-num, emp-num  $\rightarrow$  Job class

### ③ Partial Dependencies

An attribute B is functionally dependent on an attribute X which is a subset of the <sup>primary</sup> key A

★ Primary key of table is Proj-Num, Emp-Num

★ But Proj-Name is dependent on Proj-Num which is a subset of primary key

### ④ Transitive Dependency:-

Transitive dependency is the <sup>dependency</sup> of one non-prime attribute on another non-prime attribute

- Prime attribute of a relation are ~~are~~ is the 'A' attribute which is the ~~main~~ member of the candidate key of R

- Non-prime attribute of a relation are those attributes which are not a member of any candidate key of R.

2/04/19

### ⇒ Objectives

- 1) Each table represent a single subject for example a course table will contain the data that is directly related to the course
- 2) No data item unnecessarily stored in more than one table this ensures that the data is updated at one place only.
- 3) All the non-prime attribute in a table are dependent on the primary key this ensures that



delta is uniquely identifiable by the primary key  
 w/ each table is free from insertion, updation,  
 and deletion anomalies

⇒ Conversion to First Normal Form (1NF)

Step 1 - Eliminate the repeating group.

Each cell in a table has a single value and  
 there are no repeating groups.

each cell contains appropriate.

Step 2 - Identify the primary key.

eg. Proj-Num, Emp-Num

Proj-Num	Proj-Name	Emp-Numb	Emp-Name	Job-class	Chg-Hour	Hours-worked
15	Manufac	101		Manager	500	
		102		Clerk	250	
		105		Pr. Manager	600	
		110		PA	300	
		103				
16	HR	103				
		105				
		110				
		114				
18	Finance	103				
		114				
		110				



Step-3 - Identify the Dependencies

- (i) Proj. Num → Proj. Num
- (ii) Emp. Num → Emp. Name, Job-class, chrg. hours, hours
- (iii) Job. class → Chg. Hour

The first Normal form describes the tabular format

- 1) All the key attributes are identified
- 2) There are no repeating groups in the table in other words each cell in the table contains one and only one value and not the set of value.
- 3) All the attributes are dependent on the primary key.

⇒ Anomalies - (in ~~1NF~~ 1NF)

1) Insertion -

It is difficult to insert the employee details who are not yet assigned a project.

2) Update -

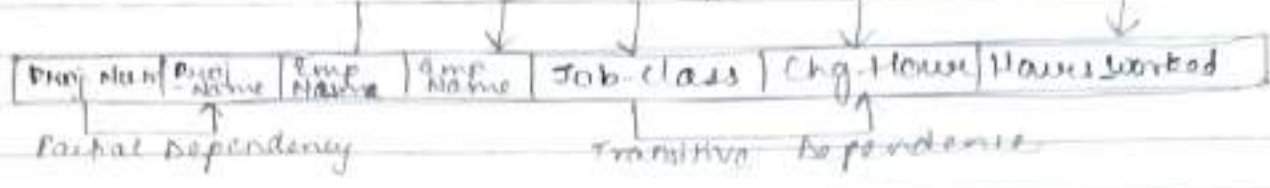
It is difficult to update employee details who are working on multiple project.

3) Deletion

If an employee working in the a project leaves the company and he is the only employee working in that project, deleting the details of the employee will also delete the project details.

NOTE - All these anomalies arise due to the existence of partial dependency.

### Partial Dependence



1NF (Proj-Num, Emp-Num, Proj-Name, Emp-Name, Job-class, Chg-hours, Hours-worked)

### Partial Dependency

Proj-Num → Proj-Name

Emp-Num → Emp-Name, Job-class, Chg-Hour, Hours-worked

### Transitive Dependency

Job-class → Chg-Hour

02/04/19

⇒ Conversion to Second Normal Form (2NF)

→ It is done when 1NF has a composite primary key.

→ If 1NF has a single attribute primary key then the table is automatically 2NF

### Steps to convert

Step 1 - Write each key component on a separate line and the original (composite key) on last line.

Proj-Num

Empl-Num

Proj-Num, Empl-Num

Each component will become a key in the new table

The original table will be divided into 3 tables -

PROJECT (Proj-Num)

Employee (Emp-Num)

Assignment (Proj-Num, Emp-Num)

Step-2 - Assign the corresponding depending attributes

PROJECT (Proj-Num, Proj-Name)

Employee (Emp-Num, Emp-Name, Job-class, chg-Hour)

Assignment (Proj-Num, Emp-Num, Hours)

2NF

Table 1: PROJECT

Proj-Num	Proj-Name
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↑ Direct relationship

Table 2: Employee

Empl-Num	emp-Name	Job-class	chg-Hour
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↑ Transitive Dependency

Table 3: Assignment

Proj-Num	Emp-Num	Assign-Hour
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⇒ Anomalies :

① Since the 2NF has transitive dependencies which will generate anomalies

② For eg:- If the change per hour changes for a job classification held by many employees then that change must be made for each of that employee.

③ If we forget to update some of the employee records then different employee with same job description will generate different hourly changes.



NOTE

A table is in 3NF when

- ① It is in 2NF and
- ② it includes no partial dependencies i.e. that no attribute is dependent on a portion of the primary key.

⇒ Conversion to 3NF:

- Step-1

Identify the determinant attributes.

- A Determinant is any attribute whose value determines other values within the a row.
- If we have 3 transitive dependencies then we will have 3 different determinants.
- In our example Job-class is the determinant.

- Step 2

Identify the dependent attributes on each determinant.

eg - Job-class → chg-hour.

- Step-3

Remove the dependent attributes from the original table

- Step 4

Create new table with determinant and its dependent attribute as a primary key.

eg - JOB (Job-class, chg-hour)

Table 1: Project

Proj-Num	Proj-Name
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Table 2: Employee

Emp-Num	Emp-Name	Job-class
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Table 3: Job

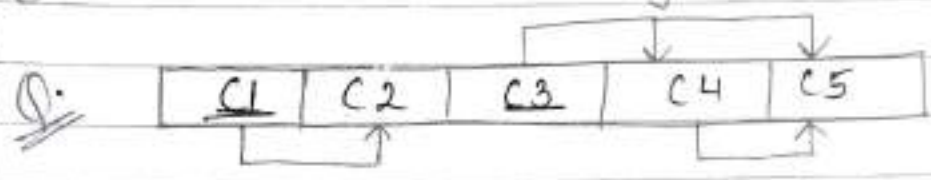
Job-class	chg-hour
-----------	----------

Table 4: Assignment

Proj-Num	Emp-Num	Assign-Hours
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A table is in 3NF -

- ① It is in 2NF
- ② It does not contain any transitive dependencies



Underlined - Primary Key - C1, C3  
 Underlined - Prime attribute  
 Non-underlined - Non-prime attribute

2NF

- Table 1 (C1, C2) // C1 is identifying C2
- Table 2 (C3, C4, C5)
- Table 3 (C1, C3) // no attribute simply write

3NF

- Table 4 (C4, C5)
- Table 2 (C3, C4)



Invoice

Inv-Num	Prod-Num	Sale-Date	Prod-Label	Vend-Code	Vend-Name	Quant-Sold	Prod-Price
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Primary key - Inv-Num, Prod-Num

Normal form - ~~2NF~~ 3NF

Dependencies - Prod-Num → Prod-Label, Prod-Price

Inv-Num → Sale-Date

Vend-Code → Vend-Name

Prod-Num, Inv-Num → Quant-Sold

2NF

Product (Prod-Num)

SELL (Inv-Num, Prod-Num)

Step-2 :- Invoice (Inv-Num, Sale-Date)

Prod (Prod-Num, Prod-Label, Prod-Price, Vend-Code, Vend-Name)

SELL (Inv-Num, Prod-Num, Qty-Sold)

3NF

Step-1:- Identify determinant

(Determinant are non-prime attribute which identify row)

In our eg:-

Vend-Code is determinant

Step-2:- Identify dependent attribute

Vend-Code → Vend-Name

Step 3. Remove dep. attribute from original table

"Product" table remove Vend-Name

- Invoice (Inv. Name, Sale. Date)
- Prod (Prod. Num, Prod. Label, Prod. Price, Vond. Code)
- Vendor (Vond. Code, Vond. Name)
- Sell (Inv. Num, Prod. Num, Qty. Sold)

Q. What are the 3 possibilities which exist when we delete a tuple that is a target of a foreign key reference?

Ans: Employee

Department

f.k.

Emp. Num | Emp. Name | D.No

D. Num | D. Name

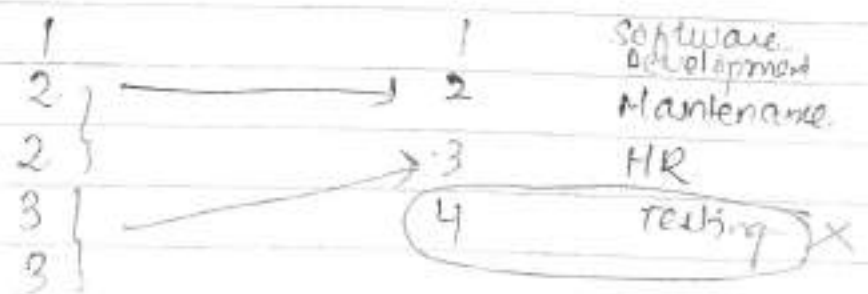
All the tuples that contain reference to the deleted tuple should also be deleted.

This is referred as Cascading Deletion.

Emp. Num	Emp. Name	D.No	D. Num	D. Name
		4	4	Testing — X
		4		
		4		

} all removed

II Only tuples that are not ~~preferred~~ referenced by another tuple can be deleted. In other words, ~~the~~ a tuple referenced by other tuple in a database cannot be deleted.



and  
 (III) The tuple is deleted ~~but~~ the F.K. attributes of all the referencing tuple are said to null.

✓

NULL

2 }

2 }

3 }

3 }

✓

1 software-development

2 Maintenance

3 HE

4 Testing